

Introduction to Quantum Information- III



General Ideas

The Goal: Define and Calculate the following:

Classical Channel

Quantum Channel

One Capacity

Four different capacities

یک کد گذاری ساده برای تصحیح خطا

x $P(y|x)$ y

0 \longrightarrow 000

1 \longrightarrow 111

نرخ مخابره

$$R = \frac{1}{3}$$

This is a very simple **letter** code.

Block Codes

010010010010001001001111010010100101001

k

010010010001011011011000100100111101000110101101110100101001

n

$M_k \longrightarrow X^{(n)}$

Encoding



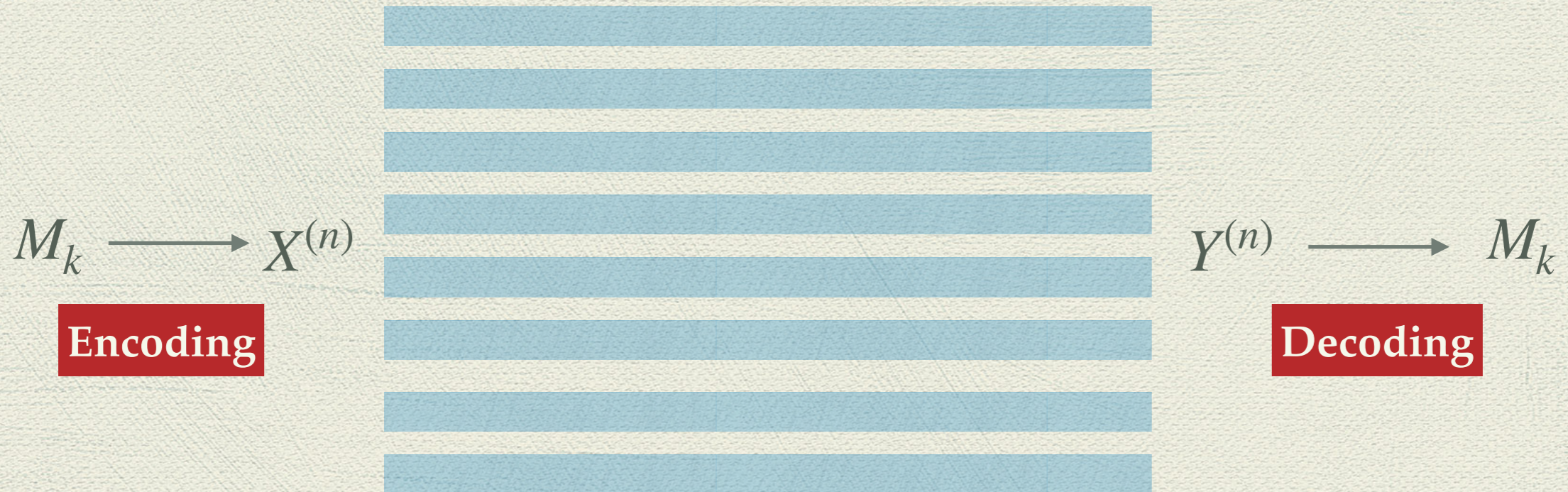
$Y^{(n)} \longrightarrow M_k$

Decoding

$$R(X) = \lim_{n \rightarrow \infty} \frac{k}{n}$$

$$P_{error} \longrightarrow 0$$

تعریف: ظرفیت کانال کلاسیک



$$C = \underset{X}{\text{Max}} R(X)$$

می بایست روی تمام منابع ها بیشینه نرخ را حساب کرد.

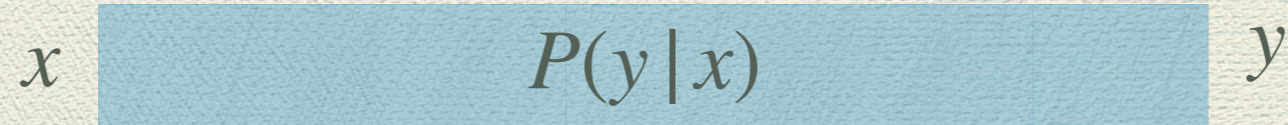
محاسبه: ظرفیت کانال کلاسیک

$$x \quad P(y|x) \quad y$$

$$I(X : Y) = H(X) + H(Y) - H(X, Y)$$

$$C = \underset{P(x)}{\text{Max}} I(X : Y)$$

محاسبه: ظرفیت کانال کلاسیک

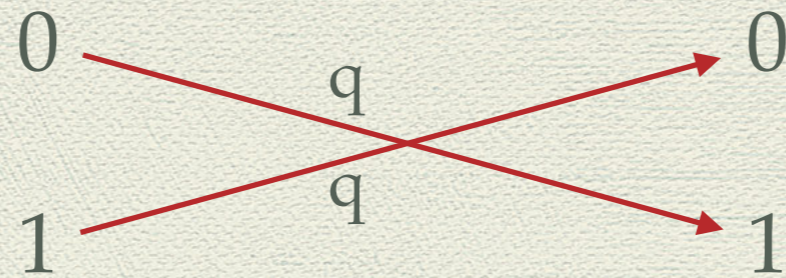


$$P(x, y) = P(y|x)P(x)$$

$$H(X) = - \sum_x P(x) \log P(x) \quad H(Y) = - \sum_y P(y) \log P(y)$$

$$H(X, Y) = - \sum_{x,y} P(x, y) \log P(x, y)$$

مثال: محاسبه ظرفیت یک کانال متقارن



$$I(X : Y) = q \log q + (1 - q) \log(1 - q) + [H(\lambda)]$$

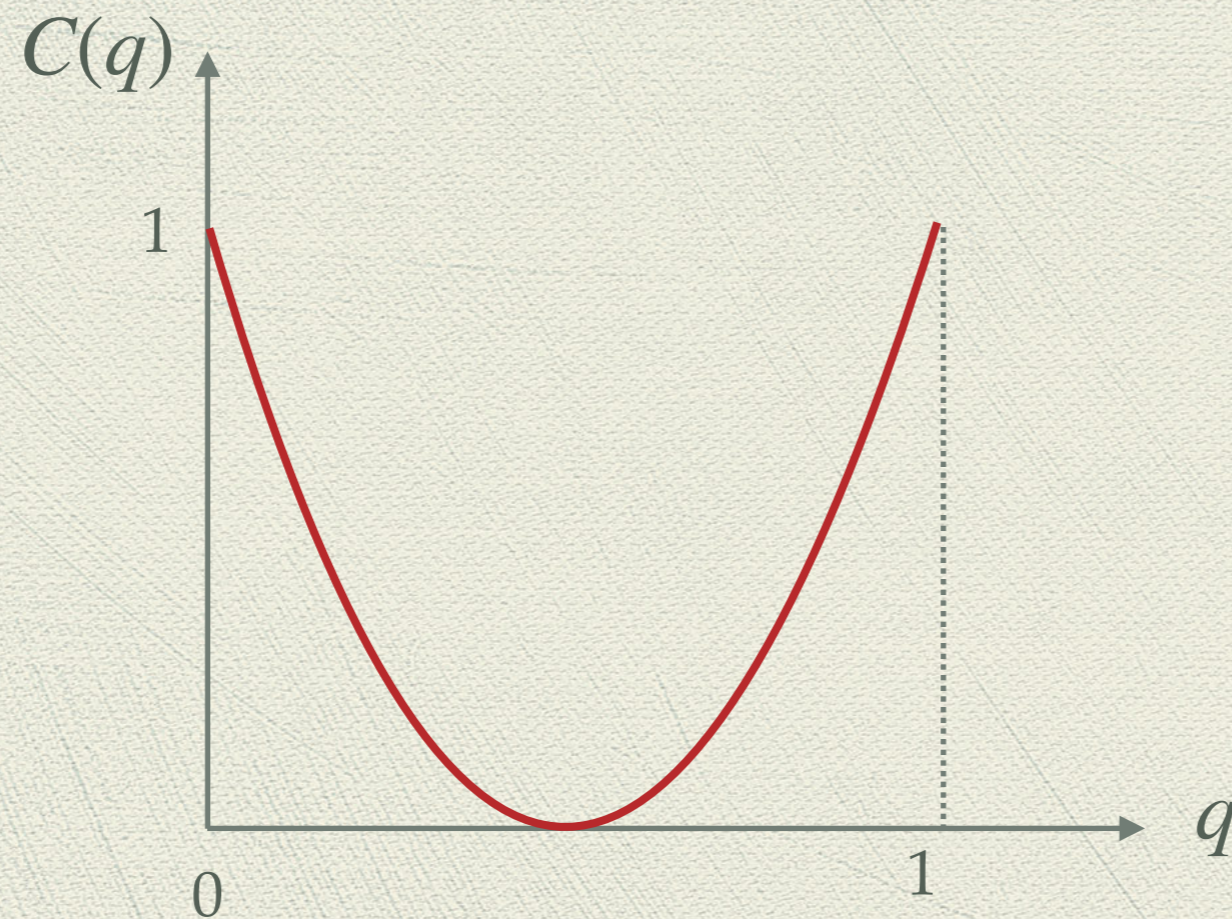
$$\lambda = p + q - 2pq$$

$$p + q - 2pq = \frac{1}{2} \longrightarrow p = \frac{1}{2}$$

بیشینه کردن

نتیجه نهایی محاسبه

$$C(q) = 1 + q \log q + (1 - q) \log(1 - q)$$



ظرفیت های کانال های کوانتومی



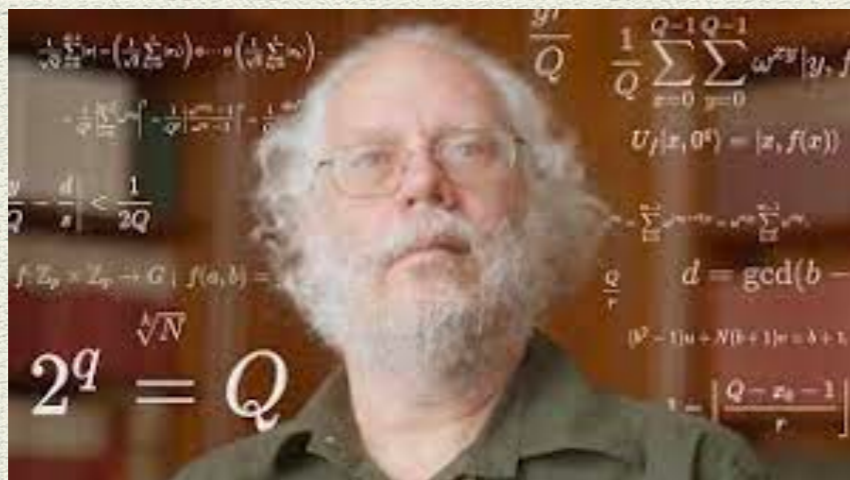
Holevo



Shumacher



Westmoreland



Peter Shor



Charles Bennett

یک- برای یک کانال کوانتومی انواعی از
ظرفیت تعریف می شود.

دو- تعریف این ظرفیت ها یک چیز است و تقلیل این تعریف ها به یک رابطه بسته یک چیز دیگر.

سه -حتی محاسبه آن رابطه بسته هم آسان نیست
(به دلیل بیشینه کردن روی یک فضای پیوسته)

چهار- برخلاف ظرفیت کلاسیک این ظرفیت ها
عموما جمع پذیر نیستند (به دلیل درهم تنیدگی).

Q1-Classical Capacity of Quantum Channels

... have served to increase the perceived value of the reward and devalue the effort. What tends to happen in the English Cup in Russia, his attitude was calmer

Definition: Classical Capacity of Quantum Channels

M=Classical message of k-bits

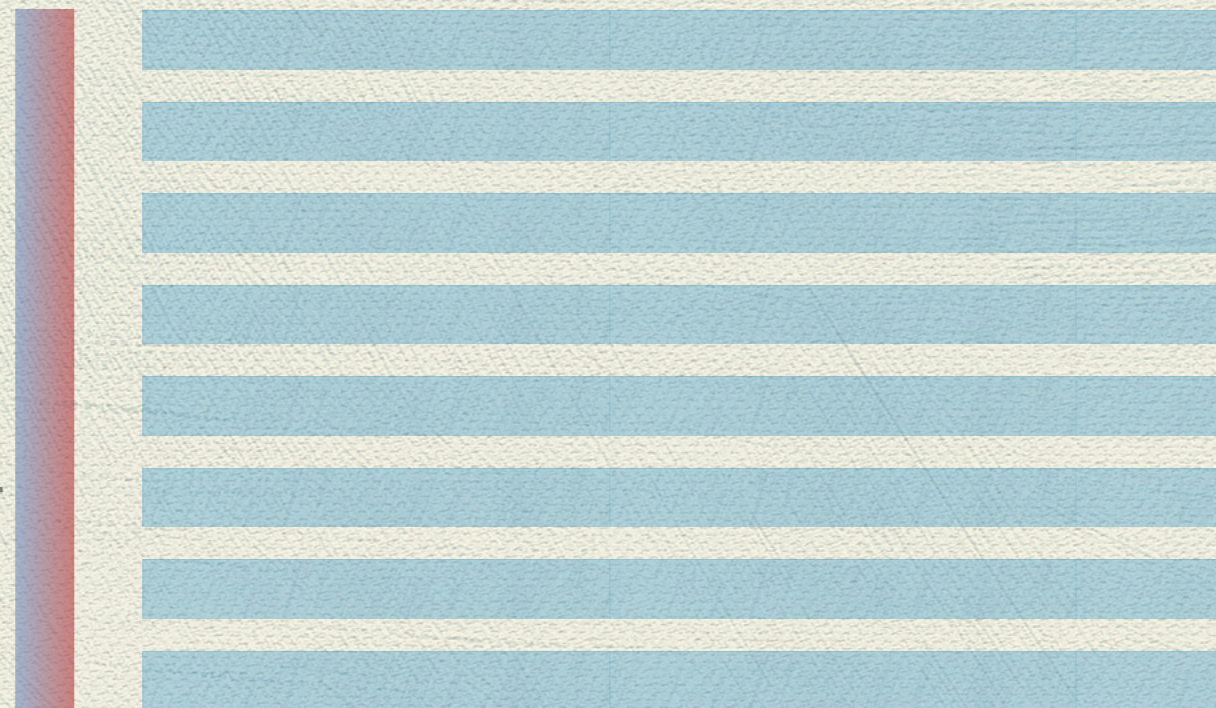


The states can be entangled.

Measurements can be entangled.

$|\psi^n(M)\rangle$ = Quantum state of n-qubits

$|\psi^n(M)\rangle$

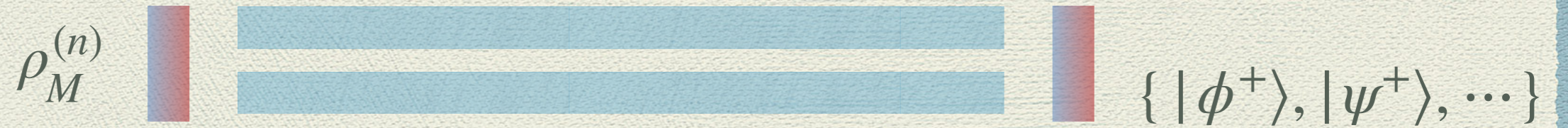


$\{E_1, E_2, \dots\}$

این حالت ها بر هم عمودند.

این حالت ها ممکن است بر هم عمود نباشند
و عموماً هم آمیخته هستند.

Example:



$$00 \longrightarrow |\phi^+\rangle$$

$$01 \longrightarrow |\psi^+\rangle$$

$$10 \longrightarrow |\psi^-\rangle$$

$$11 \longrightarrow |\phi^-\rangle$$

Quantum Channel

$$\rho_{00}$$

$$\rho_{01}$$

$$\rho_{10}$$

$$\rho_{11}$$

$$P_{error}(00) = 1 - \langle \phi^+ | \rho_{00} | \phi^+ \rangle$$

$$P_{error}(01) = 1 - \langle \psi^+ | \rho_{01} | \psi^+ \rangle$$

.....

.....

$$P_{error} = \frac{P_{error}(00) + P_{error}(01) + \dots}{4}$$



$$P_{error}(M) = 1 - Tr\left(E_M \Phi^{\otimes n}(\rho_M^{(n)})\right)$$

$$\bar{P}_{error} = \frac{\sum_M 1 - Tr\left(E_M \Phi^{\otimes n}(\rho_M^{(n)})\right)}{|M|}$$

Achievable rate for the source X

$$R(X) = \lim_{n \rightarrow \infty} \frac{\log |M|}{n} = \lim_{n \rightarrow \infty} \frac{k}{n}$$

when $P_{error} \rightarrow 0$

تعریف ظرفیت کلاسیک یک کانال کوانتومی

$$C := \underset{X}{\text{Max}} R(X)$$

ماکزیمم گیری روی تمام منابع ها صورت می گیرد.

این همان تعریف کلی و دقیق ظرفیت است و هنوز فرمول بسته نیست.

فرمول بسته برای ظرفیت کلاسیک کانال کوانتومی؟

فقط در یک حالت وجود دارد.

$C^{(1)}$

$\rho_1 \otimes \rho_2 \otimes \dots \otimes \rho_n$



$\{E_M, \dots\}$

وقتی حالت های ورودی به صورت ضربی باشند: $\rho_1 \otimes \rho_2 \otimes \dots \otimes \rho_n$

فرم بسته برای ظرفیت کلاسیک کانال کوانتومی در یک حالت خاص: $C^{(1)}$

Holevo-Schumacher-Westmoreland



$$\rho_1 \otimes \rho_2 \otimes \dots \otimes \rho_M$$



$$\{E_M, \dots\}$$

Holevo-Schumacher-Westmoreland Or Product State Capacity

$$C^{(1)} = \chi^*(\Phi)$$

$$\chi^*(\Phi) = \underset{\{p_i, \rho_i\}}{\text{Max}} \chi(\{p_i, \rho_i\})$$

$$\chi(\Phi) = S\left[\sum_i p_i \Phi(\rho_i)\right] - \sum_i p_i S(\Phi(\rho_i))$$

Example 1: $C^{(1)}$ Capacity of **Depolarizing** Channel

$$\Phi(\rho) = (1 - p)\rho + \frac{p}{2}I$$

We take an ensemble $\{p_i, \rho_i\}$ and calculate

$$\chi(\Phi) = S\left[\sum_i p_i \Phi(\rho_i)\right] - \sum_i p_i S(\Phi(\rho_i))$$

$$\sum_i p_i \Phi(\rho_i) = \Phi\left(\sum_i p_i \rho_i\right) = \Phi(\rho) \qquad \Phi(\rho_i) = (1 - p)\rho_i + \frac{p}{2}I$$

$$\chi(\{p_i, \rho_i\}) = S\left[(1-p)\rho + p\frac{I}{2}\right] - \sum_i p_i S\left[(1-p)\rho_i + \frac{p}{2}I\right]$$

The input ensemble can be taken to be pure,
Thm. 13.3.2. Mark Wilde ([arxiv.org.1106.1445.](https://arxiv.org/abs/1106.1445)).

To maximize the first term, we take $\rho = \frac{I}{2}$.

For the second term:

$$\text{Since } S(U\sigma U^\dagger) = S(\sigma) \quad \forall U, \sigma$$

$$S\left[(1-p)|\psi_i\rangle\langle\psi_i| + p\frac{I}{2}\right] = S\left[(1-p)|+\rangle\langle+| + p\frac{I}{2}\right]$$

**Therefore it is important to know
the minimum output entropy states.**

$$\chi(\{p_i, \rho_i\}) = S\left[(1-p)\rho + p\frac{I}{2}\right] - \sum_i p_i S\left[(1-p)\rho_i + \frac{p}{2}I\right]$$

$$= 1 - S\left[(1-p)\left|+\right\rangle\langle+| + \frac{p}{2}I\right]$$

$$C^{(1)}(p) = 1 + \frac{p}{2} \log \frac{p}{2} + \left(1 - \frac{p}{2}\right) \log \left(1 - \frac{p}{2}\right)$$

$$C^{(1)}(0) = 1$$

$$C^{(1)}(1) = 0$$

Example 2: $C^{(1)}$ Capacity of **Bit-flip** Channel

$$\Phi(\rho) = (1 - p)\rho + pX\rho X$$

We take $\rho = \frac{I}{2}$ to maximize the first term.

$$\chi(\{p_i, \rho_i\}) = S[(1 - p)\rho + pX\rho X] - \sum_i p_i S[(1 - p)\rho_i + pX\rho_i X]$$

Intuition: $\{p_i, |\psi_i\rangle\} = \{1/2, |+\rangle; 1/2, |-\rangle\}$


Invariant States

$$\chi(\{p_i, \rho_i\}) = S[(1-p)\rho + pX\rho X] - \sum_i p_i S[(1-p)\rho_i + pX\rho_i X]$$

Intuition: $\{p_i, |\psi_i\rangle\} = \{1/2, |+\rangle; 1/2, |-\rangle\}$

$$\chi = S\left(\frac{I}{2}\right) - \frac{1}{2}S(|+\rangle\langle +|) - \frac{1}{2}S(|-\rangle\langle -|) = 1$$

C=1 as it should be.

Example 3: $C^{(1)}$ Capacity of **Amplitude-Damping** Channel

$$\Phi(\rho) = A_0\rho A_0^\dagger + A_1\rho A_1^\dagger$$

$$A_0 = \begin{pmatrix} 1 & 0 \\ 0 & \sqrt{1-\gamma} \end{pmatrix} \quad A_1 = \begin{pmatrix} 0 & \sqrt{\gamma} \\ 0 & 0 \end{pmatrix}$$

Intuition: $\{p_i, |\psi_i\rangle\} = \{p, |0\rangle; 1-p, |\psi\rangle\}$

Intuition: $\{p_i, |\psi_i\rangle\} = \{p, |0\rangle; 1 - p, |\psi\rangle\}$

$$|\psi\rangle = \begin{pmatrix} a \\ b \end{pmatrix}$$

$$|0\rangle\langle 0| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$|\psi\rangle\langle\psi| = \begin{pmatrix} |a|^2 & a\bar{b} \\ \bar{a}b & |b|^2 \end{pmatrix}$$

$$\rho = \begin{pmatrix} p + (1 - p)|a|^2 & (1 - p)a\bar{b} \\ (1 - p)\bar{a}b & (1 - p)|b|^2 \end{pmatrix}$$

$$\Phi(|0\rangle\langle 0|) = |0\rangle\langle 0|$$

$$\Phi(|\psi\rangle\langle\psi|) = \begin{pmatrix} |a|^2 + \gamma|b|^2 & a\bar{b} \\ \bar{a}b & (1 - \gamma)|b|^2 \end{pmatrix}$$

$$\Phi(\rho) = \begin{pmatrix} p + (1 - p)|a|^2 + \gamma(1 - p)|b|^2 & (1 - p)a\bar{b} \\ (1 - p)\bar{a}b & (1 - \gamma)(1 - p)|b|^2 \end{pmatrix}$$

$$\chi_\gamma(p, a) = S[\Phi(\rho)] - (1 - p)S[\Phi(|\psi\rangle\langle\psi|)]$$

$$\chi(p, a) = S[\Phi(\rho)] - (1 - p)S[\Phi(|\psi\rangle\langle\psi|)]$$

Using $S(\sigma) = S(U\sigma U^\dagger)$ we can take a and b to be real.

$$a = \cos \theta$$

$$b = \sin \theta$$

$$\chi(p, \theta) = S[\Phi(\rho)] - (1 - p)S[\Phi(|\psi\rangle\langle\psi|)]$$

χ^* and hence C can be calculated only numerically.



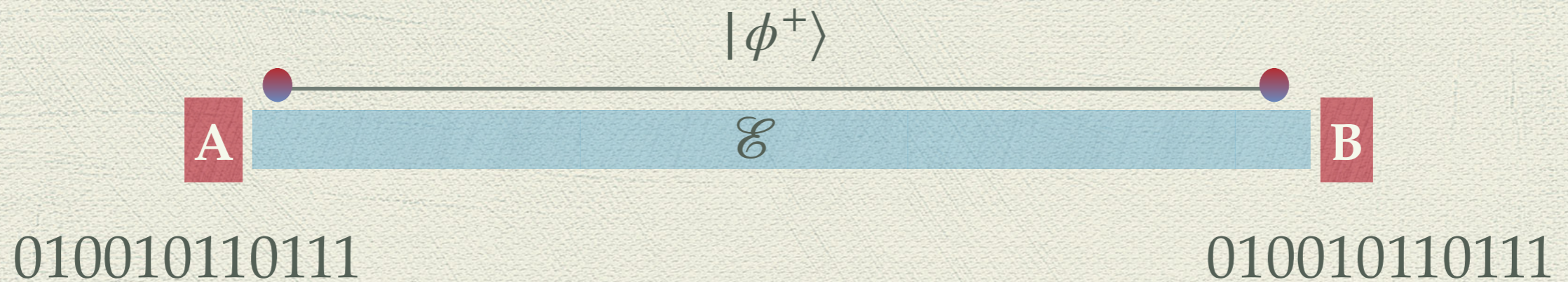
Q2-Entanglement-Assisted Classical Capacity of Quantum Channels

... have served to increase the number of children born to women who are thinking would then be: "It must be serious."
"What tends to happen in the English Premier League?"
emphasise the reward and devalue the...
2018, when Gareth Southgate managed England to win the World Cup in Russia, his attitude was calmer and...

Q2: Entanglement-Assisted Classical Capacity of Quantum Channels

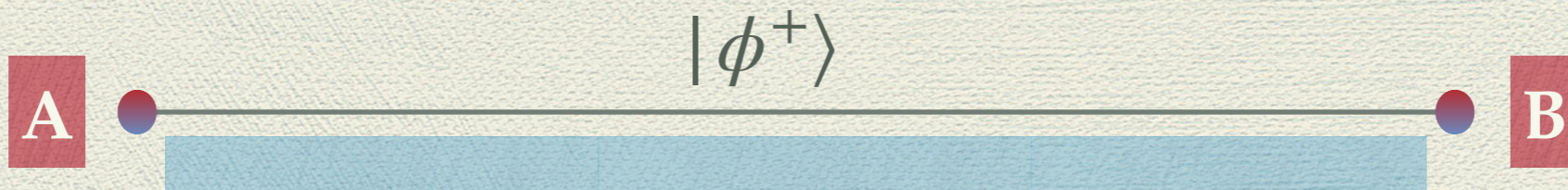
بین فرستنده و گیرنده به دلخواه حالت های درهم تنیده ماکزیمال وجود دارد که می توانند از آن برای ارسال پیام های کلاسیک استفاده کنند.

Definition: Entanglement-assisted Classical Capacity
of Quantum Channels



Example: Dense Coding

Example 1: Dense Coding

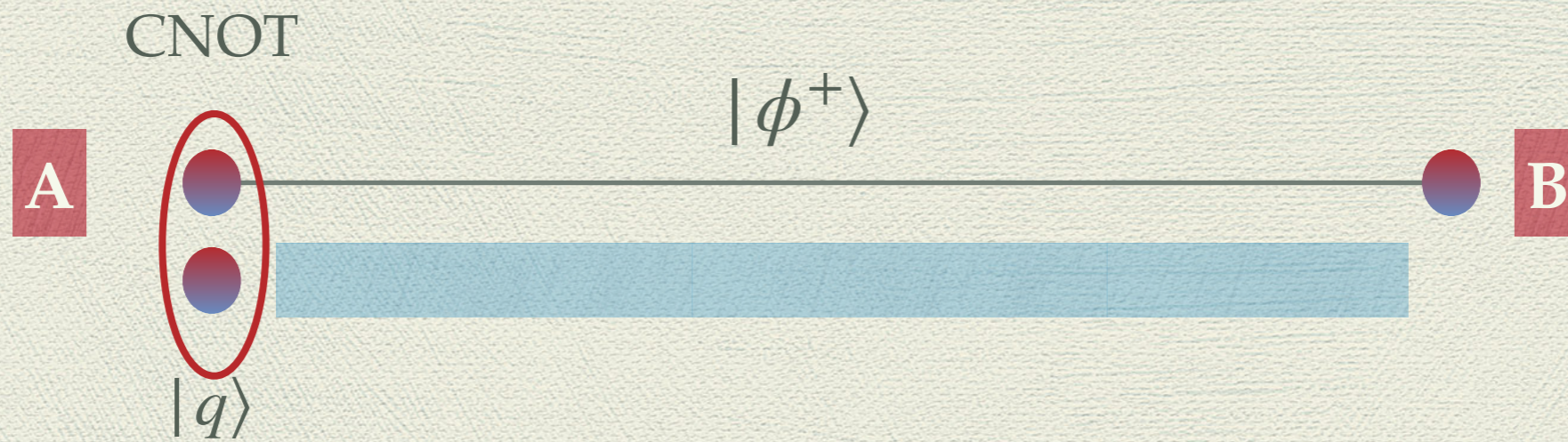


$$|\phi^+\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

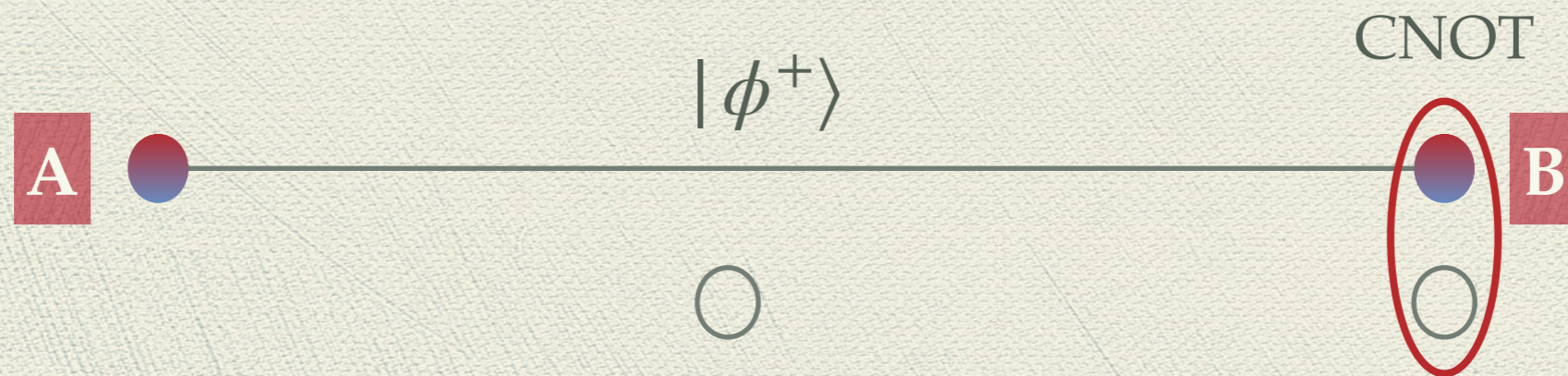
Arrows labeled I , X , Y , and Z point from the initial state to the following equations:

$$|\phi^+\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) \quad 00$$
$$|\phi^+\rangle = \frac{1}{\sqrt{2}}(|10\rangle + |01\rangle) \quad 01$$
$$|\phi^+\rangle = \frac{1}{\sqrt{2}}(|10\rangle - |01\rangle) \quad 10$$
$$|\phi^+\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle) \quad 11$$

Example 2



$$|\Psi\rangle_{ABA} = \frac{1}{\sqrt{2}}(|00q\rangle + |11\bar{q}\rangle)$$



General Definition: Entanglement-assisted Classical Capacity of Quantum Channels

$$|\phi^+\rangle^{\otimes \infty}$$



$\rho_M^{(n)}$ = Quantum state of n-qubits

$$C = \text{Max}_{n \rightarrow \infty} \frac{|M|}{n} \quad ?$$

Calculation: Entanglement-Assisted Classical Capacity of Quantum Channel

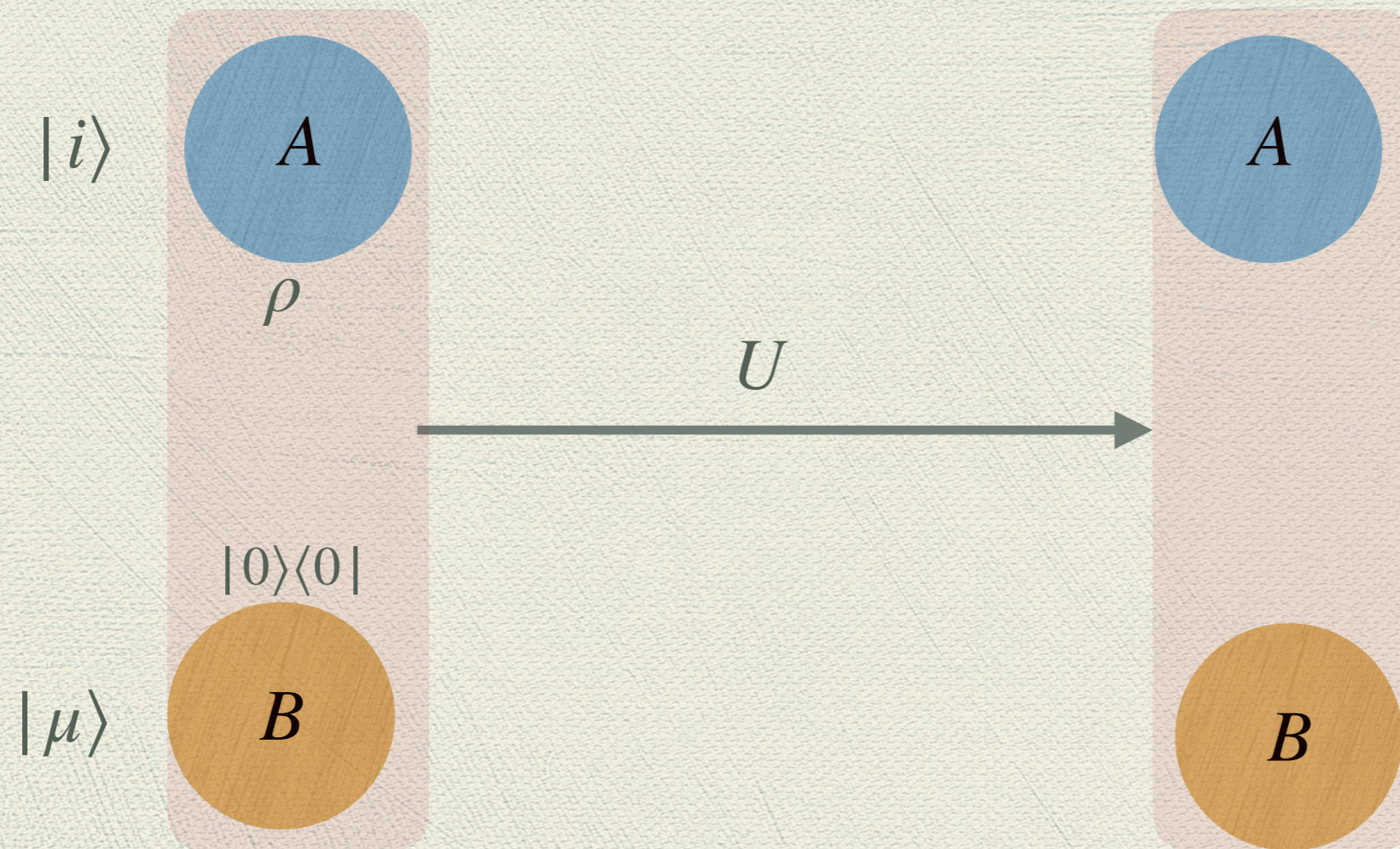
$$C_E(\Phi) = \underset{\rho}{\text{Max}} I(\rho, \Phi)$$

$$I(\rho, \Phi) = S(\rho) + S(\Phi(\rho)) - S(\Phi^c(\rho))$$

بنابراین برای این نوع ظرفیت یک فرمول بسته و ساده وجود دارد.

ولی بازهم می بایست این عبارت آخری را روی تمام ماتریس های چگالی بیشینه کرد.

Complementary Channels

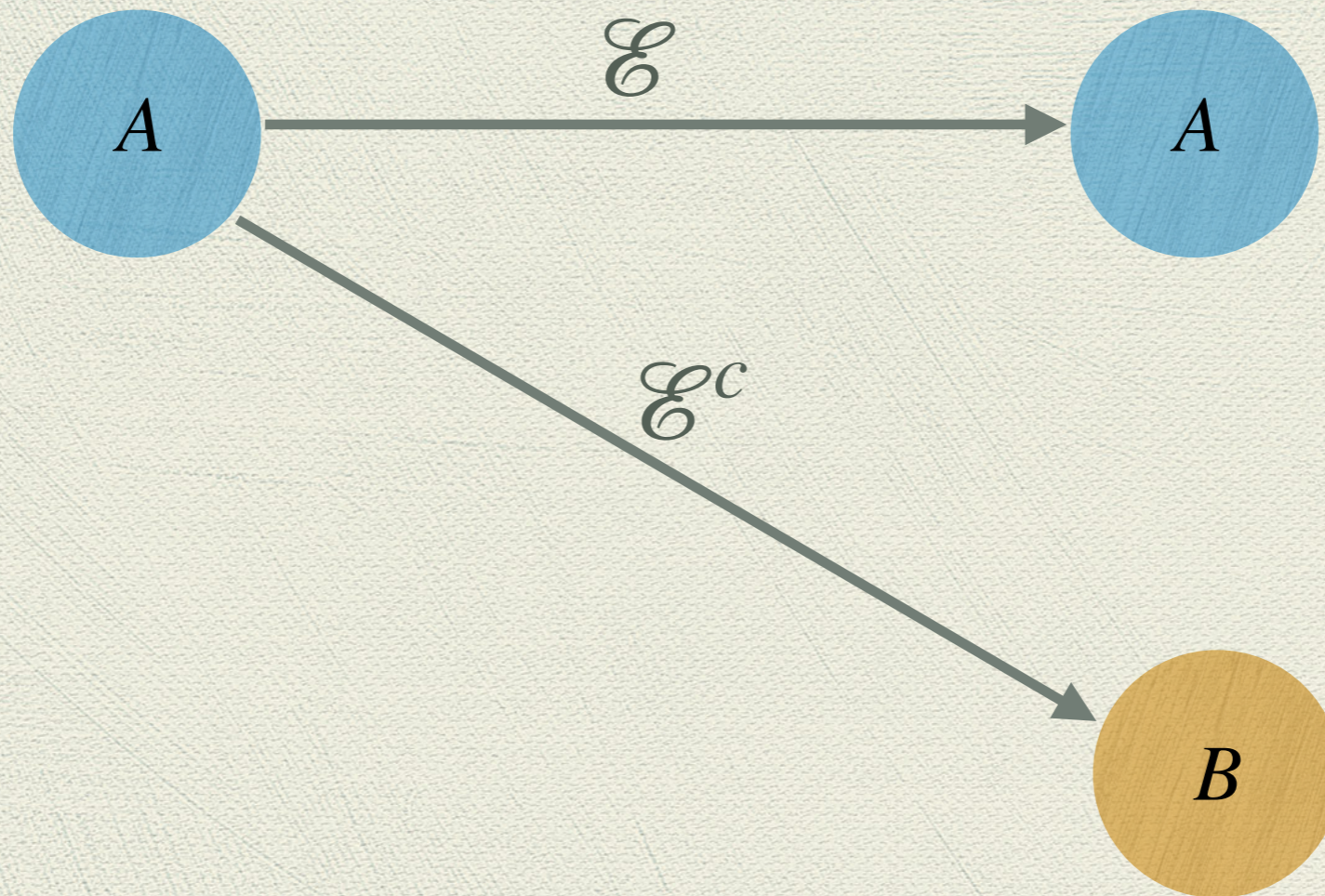


Complementary Channels

$$\mathcal{E}(\rho) = \text{Tr}_B \left[U(\rho \otimes |0\rangle\langle 0|)U^\dagger \right] \in A$$

$$\mathcal{E}^c(\rho) = \text{Tr}_A \left[U(\rho \otimes |0\rangle\langle 0|)U^\dagger \right] \in B$$

Complementary Channels



Kraus Decomposition

$$\mathcal{E}(\rho) = \sum_{\mu} A_{\mu} \rho A_{\mu}^{\dagger}$$

$$(K_i)_{\mu,j} = (A_{\mu})_{ij}$$

$$\mathcal{E}^c(\rho) = \sum_i K_i \rho K_i^{\dagger}$$

ماتریس K_1 را از سطرهای اول A ها می سازیم.

ماتریس K_2 را از سطرهای دوم A ها می سازیم.

ماتریس K_3 را از سطرهای سوم A ها می سازیم.

\mathcal{E}



A_1

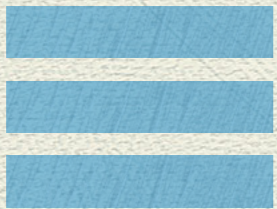


A_2



A_3

\mathcal{E}^c



K_1



K_2



K_3



K_4

Example 1: Complement of Bit-flip Channel

$$\mathcal{E}(\rho) = (1 - p)\rho + p\sigma_x\rho\sigma_x$$

$$A_0 = \begin{pmatrix} \sqrt{1-p} & 0 \\ 0 & \sqrt{1-p} \end{pmatrix} \quad A_1 = \begin{pmatrix} 0 & \sqrt{p} \\ \sqrt{p} & 0 \end{pmatrix}$$

$$K_0 = \begin{pmatrix} \sqrt{1-p} & 0 \\ 0 & \sqrt{p} \end{pmatrix} \quad K_1 = \begin{pmatrix} 0 & \sqrt{1-p} \\ \sqrt{p} & 0 \end{pmatrix}$$

$$\rho = \frac{1}{2}(I + \mathbf{r} \cdot \sigma)$$

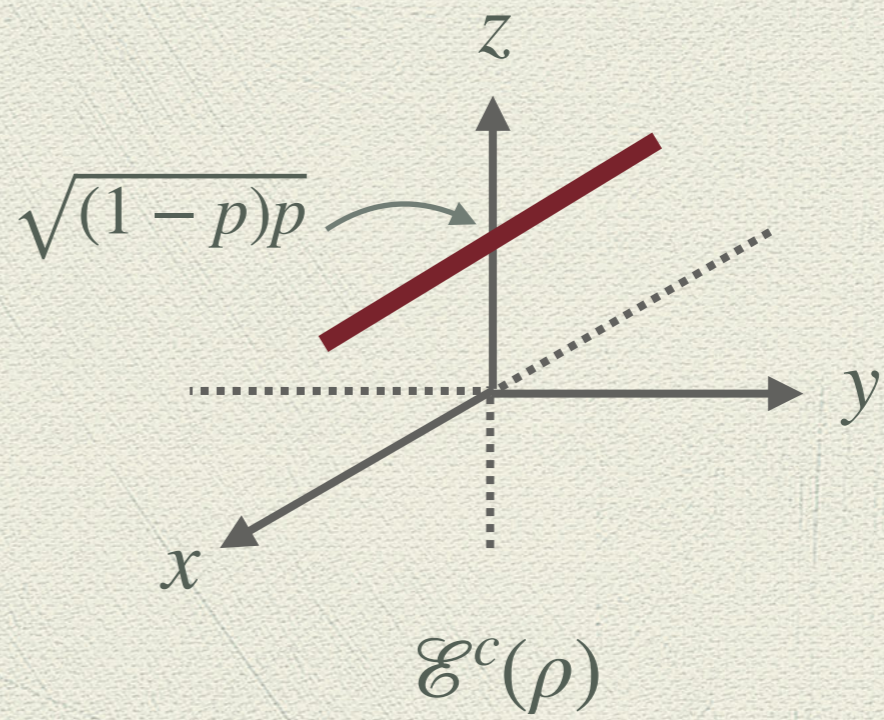
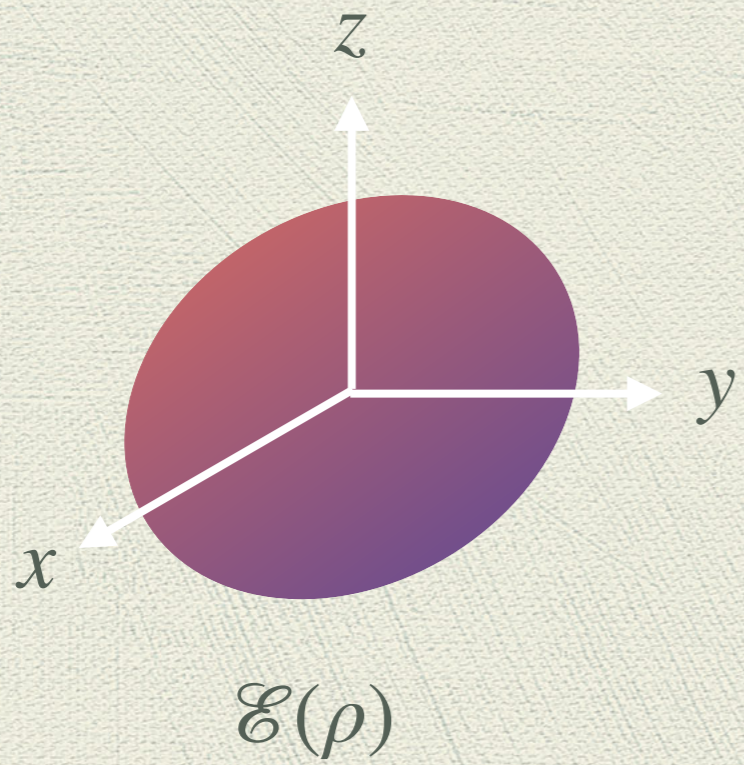
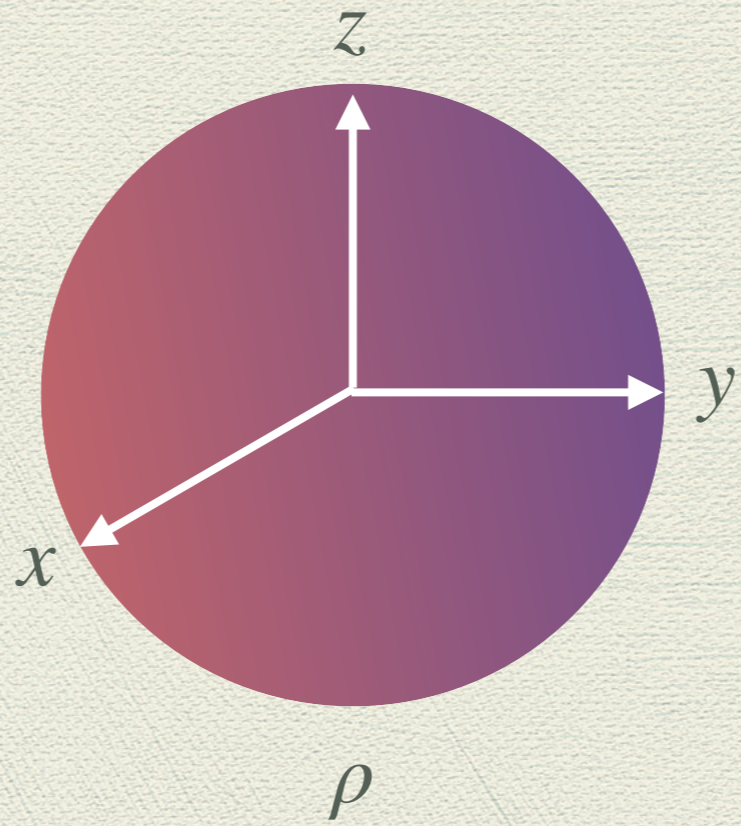
$$\mathbf{r} = (x, y, z)$$

$$\mathcal{E}(\rho) = \frac{1}{2}(I + \mathbf{r}' \cdot \sigma)$$

$$\mathbf{r}' = (x, (1 - 2p)y, (1 - 2p)z)$$

$$\mathcal{E}^c(\rho) = \frac{1}{2}(I + \mathbf{r}_c' \cdot \sigma)$$

$$\mathbf{r}_c' = (\sqrt{(1-p)p} x, 0, (1 - 2p))$$



$$I(\rho, \Phi) = S(\rho) + S(\Phi(\rho)) - S(\Phi^c(\rho))$$

Example 2: Depolarizing Channel

$$\mathcal{E}(\rho) = \left(1 - \frac{3p}{4}\right)\rho + \frac{p}{4}(\sigma_x\rho\sigma_x + \sigma_y\rho\sigma_y + \sigma_z\rho\sigma_z)$$

$$A_0 = \alpha \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$A_1 = \beta \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$A_2 = \beta \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$$

$$A_3 = \beta \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$\alpha = \sqrt{1 - \frac{3p}{4}}$$

$$\beta = \sqrt{\frac{p}{4}}$$

$$A_0 = \alpha \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$A_1 = \beta \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$A_2 = \beta \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$$

$$A_3 = \beta \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$R_1 = \begin{pmatrix} \alpha & 0 \\ 0 & \beta \\ 0 & -i\beta \\ \beta & 0 \end{pmatrix}$$

$$R_2 = \begin{pmatrix} 0 & \alpha \\ \beta & 0 \\ i\beta & 0 \\ 0 & -\beta \end{pmatrix}$$

$$\mathcal{E}^c(\rho) = \begin{pmatrix} \alpha^2 & \alpha\beta x & \alpha\beta y & \alpha\beta z \\ \alpha\beta x & \beta & -i\beta z & i\beta y \\ \alpha\beta y & i\beta z & \beta & -i\beta x \\ \alpha\beta z & -i\beta y & i\beta x & \beta \end{pmatrix}$$

Eigenvalues of $\rho : \{\frac{1}{2}(1 \pm r)\}$

Eigenvalues of $\mathcal{E}(\rho) : \{\frac{1}{2}(1 \pm (1 - p)r)\}$

Eigenvalues of $\mathcal{E}^c(\rho) : \{\beta^2(1 \pm r), \frac{1}{2}(1 - 2\beta^2 \pm \sqrt{(1 - 4\beta^2)^2 - 8(\beta^2 r^2)^2})\}$

$$I(\rho, \Phi) = S(\rho) + S(\Phi(\rho)) - S(\Phi^c(\rho))$$

Q3-Quantum Capacity of Quantum Channels

... have served to increase the ...
... thinking would then be: 'It must ...
... What tends to happen in the Em ...
... emphasise the reward and devalue ...
... 2018, when Gareth Southgate manage ...
... Cup in Russia, his attitude was cal ...

... between players and ...
... of the task (the ...
... us if I'm not even ...
... ess is that they ...
... Harkness ...
... and at the ...
... d more

Definition: Quantum Capacity
of Quantum Channels



$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

Quantum Channel Φ



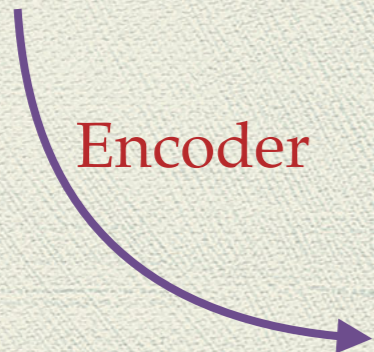
ρ

Definition: Quantum Capacity of Quantum Channels

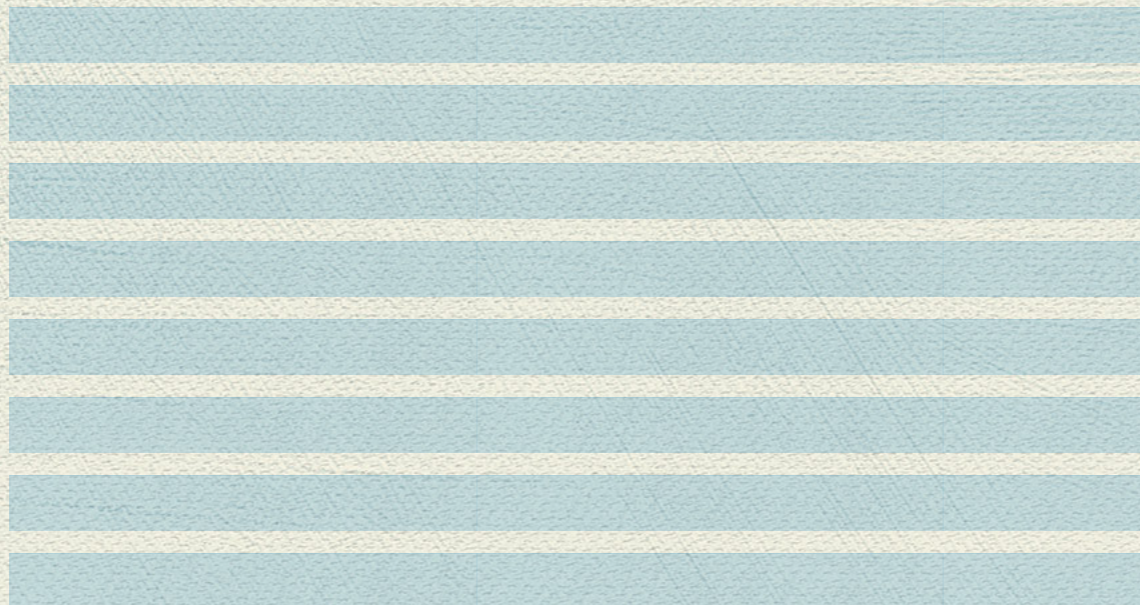
$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$



Encoder

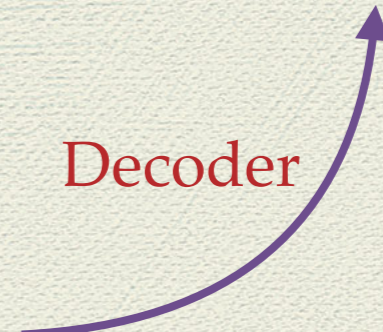


$$|\Psi\rangle_n$$



$$\Phi^{\otimes n}(|\Psi\rangle_n)$$

Decoder



$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$



Definition: Quantum Capacity of Quantum Channels

$$X = \{p_i, |\psi_i\rangle \in C^k\}$$

$$Y = \{p_i, |\psi'_i\rangle \in C^k\}$$



We want $\bar{F}(|\psi'_k\rangle, |\psi_k\rangle) \geq 1 - \epsilon$

$$R(X) = \lim_{n \rightarrow \infty} \frac{k}{n}$$

Definition: Quantum Capacity of Quantum Channels

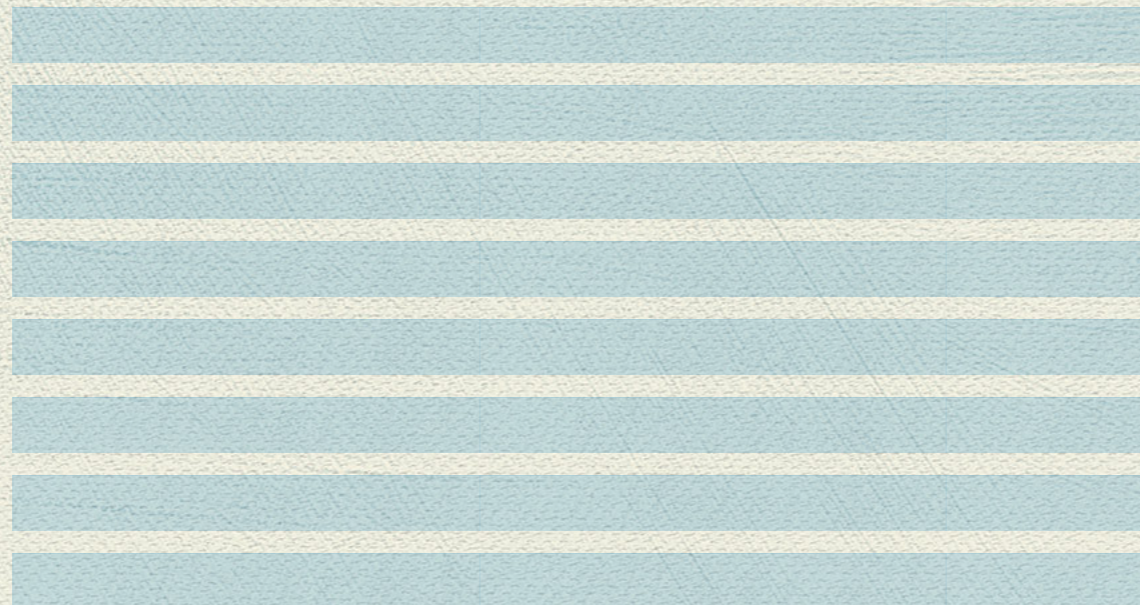
$$X = \{p_i, |\psi_i\rangle \in C^k\}$$

$$Y = \{p_i, |\psi'_i\rangle \in C^k\}$$



Encoder

$$|\Psi\rangle_n$$



Decoder

$$\Phi^{\otimes n}(|\Psi\rangle_n)$$

$$C = \underset{X}{\text{Max}} R(X)$$

Calculation: Quantum Capacity of Quantum Channels

$$J_n(\rho, \Phi) = S(\Phi^{\otimes n}(\rho)) - S((\Phi^c)^{\otimes n}(\rho))$$

$$J_n(\Phi) = \underset{\rho}{\text{Max}} J_n(\rho, \Phi^{\otimes n})$$

$$C = \lim_{n \rightarrow \infty} \frac{1}{n} J_n[\Phi]$$

Obviously this capacity cannot be calculated,

Except when the channel is degradable.

For degradable channels

$$C = J[\Phi] = \underset{\rho}{\text{Max}} S[\Phi(\rho)] - S[\Phi^c(\rho)]$$

Q4-Private Capacity of Quantum Channels

... have served to increase the ...
... thinking would then be: 'It must ...
... What tends to happen in the Em ...
... emphasise the reward and devalue ...
... 2018, when Gareth Southgate manage ...
... Cup in Russia, his attitude was cal ...

... between players and ...
... of the task (the ...
... us if I'm not even ...
... ess is that they ...
... Harkness ...
... and at the ...
... d more

M=Classical Message

M'



$$Pr(M \neq M') \leq \epsilon$$

$$\|\rho_{BE} - \rho_B \otimes \rho_E\| \leq \delta$$

Transferring classical messages to Bob without any leakage to Eve

Single Shot Capacity

$$C_p^1(\Lambda) = \max_{\{p_i, \rho_i\}} [\chi(\Lambda) - \chi(\Lambda^c)]$$

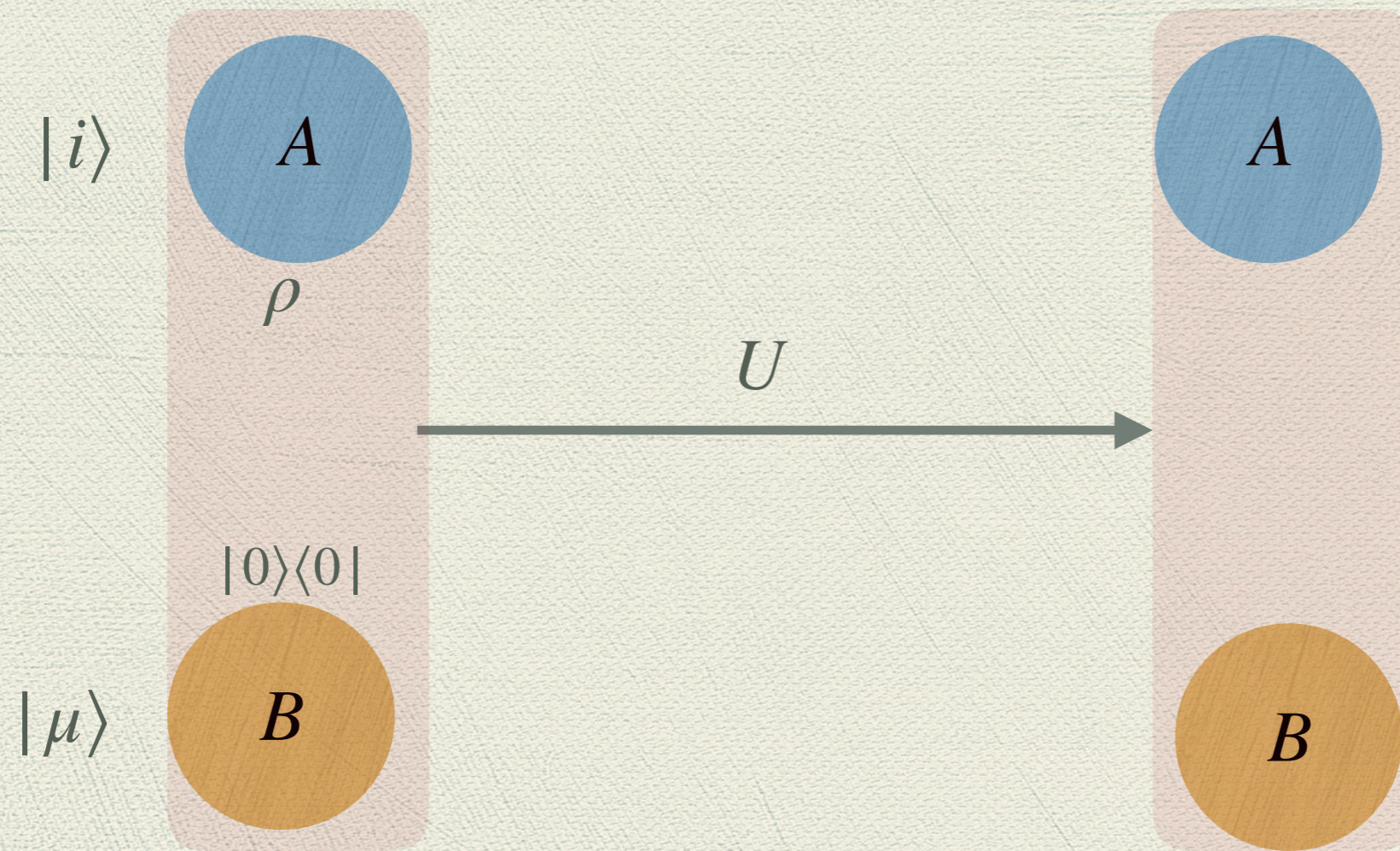
$$C_p(\Lambda) = \max_{\{p_i, \rho_i\}} \lim_{n \rightarrow \infty} [\chi(\Lambda^{\otimes n}) - \chi(\Lambda^{c \otimes n})]$$

$$C_p^1(\Lambda) \leq C_p(\Lambda)$$

$$Q(\Lambda) \leq C_p(\Lambda)$$

Degradable Channels

Complementary Channels

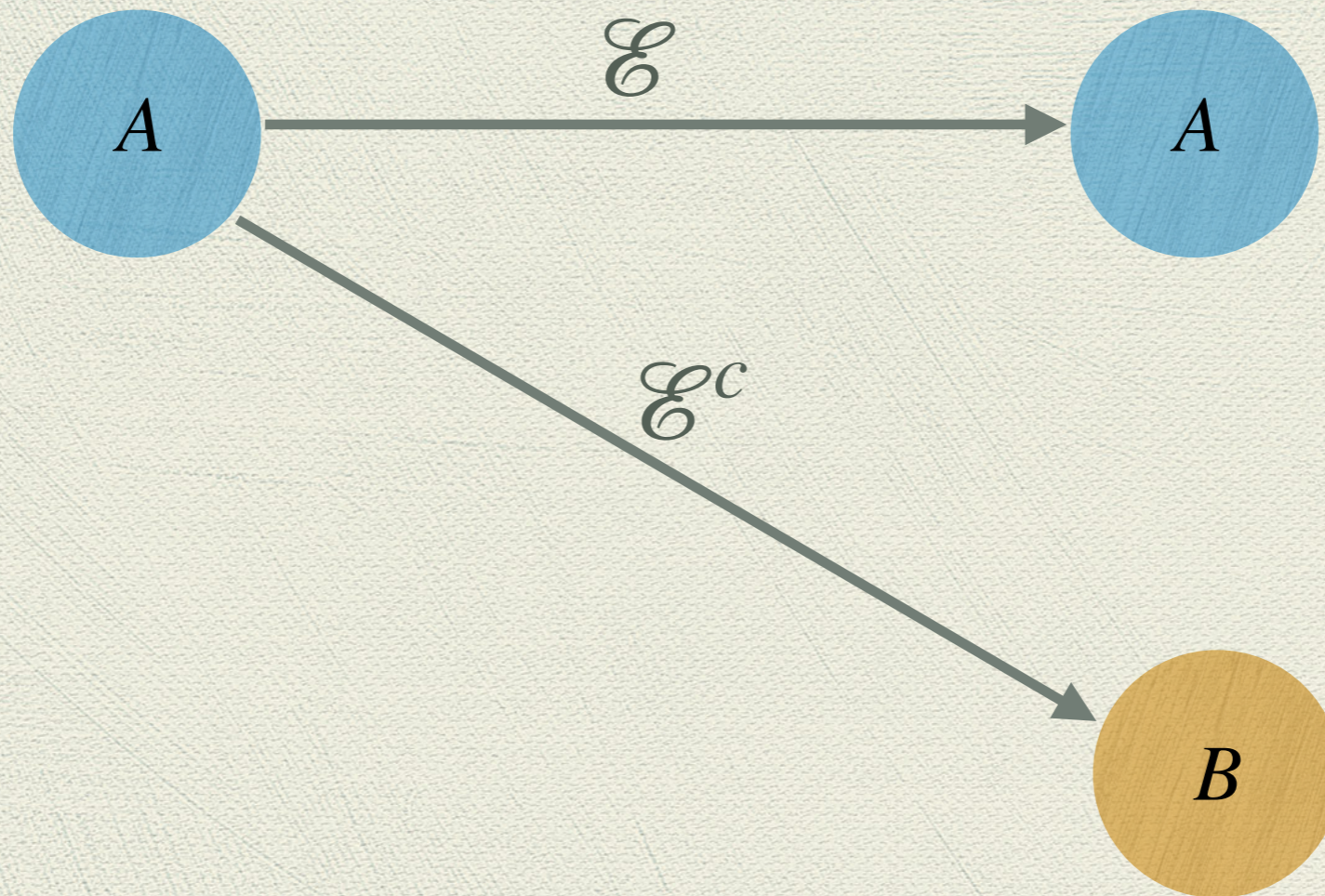


Complementary Channels

$$\mathcal{E}(\rho) = \text{Tr}_B \left[U(\rho \otimes |0\rangle\langle 0|)U^\dagger \right] \in A$$

$$\mathcal{E}^c(\rho) = \text{Tr}_A \left[U(\rho \otimes |0\rangle\langle 0|)U^\dagger \right] \in B$$

Complementary Channels



Kraus Decomposition

$$\mathcal{E}(\rho) = \sum_{\mu} A_{\mu} \rho A_{\mu}^{\dagger} \quad (A_{\mu})_{ij} = \langle i, \mu | U | j, 0 \rangle$$

$$\mathcal{E}^c(\rho) = \sum_i K_i \rho K_i^{\dagger} \quad (K_i)_{\mu,j} = \langle i, \mu | U | j, 0 \rangle$$

$$(K_i)_{\mu,j} = (A_{\mu})_{ij}$$

$$(K_i)_{\mu,j} = (A_\mu)_{ij}$$

d_A

d_A

d_B

K_i

d_A

d_A

A_μ

Example 1: Complement of Bit-flip Channel

$$(K_i)_{\mu,j} = (A_\mu)_{i,j}$$

$$\mathcal{E}(\rho) = (1-p)\rho + p\sigma_x\rho\sigma_x$$

$$A_0 = \begin{pmatrix} \sqrt{1-p} & 0 \\ 0 & \sqrt{1-p} \end{pmatrix} \quad A_1 = \begin{pmatrix} 0 & \sqrt{p} \\ \sqrt{p} & 0 \end{pmatrix}$$

$$K_0 = \begin{pmatrix} \sqrt{1-p} & 0 \\ 0 & \sqrt{p} \end{pmatrix} \quad K_1 = \begin{pmatrix} 0 & \sqrt{1-p} \\ \sqrt{p} & 0 \end{pmatrix}$$

Example 2: Complement of Pauli Channel

$$\mathcal{E}(\rho) = p_0\rho + p_1\sigma_x\rho\sigma_x + p_2\sigma_y\rho\sigma_y + p_3\sigma_z\rho\sigma_z$$

$$K_0 = \begin{pmatrix} \sqrt{p_0} & 0 \\ 0 & \sqrt{p_1} \\ 0 & -i\sqrt{p_2} \\ \sqrt{p_3} & 0 \end{pmatrix}$$

$$K_1 = \begin{pmatrix} 0 & \sqrt{p_0} \\ \sqrt{p_1} & 0 \\ i\sqrt{p_2} & 0 \\ 0 & -\sqrt{p_3} \end{pmatrix}$$

More Examples

$$\mathcal{E}_0(\rho) = \text{Tr}(\rho) \frac{I}{d}$$

$$\tilde{\mathcal{E}}_0(\rho) = \frac{I}{d} (\rho \otimes I)$$

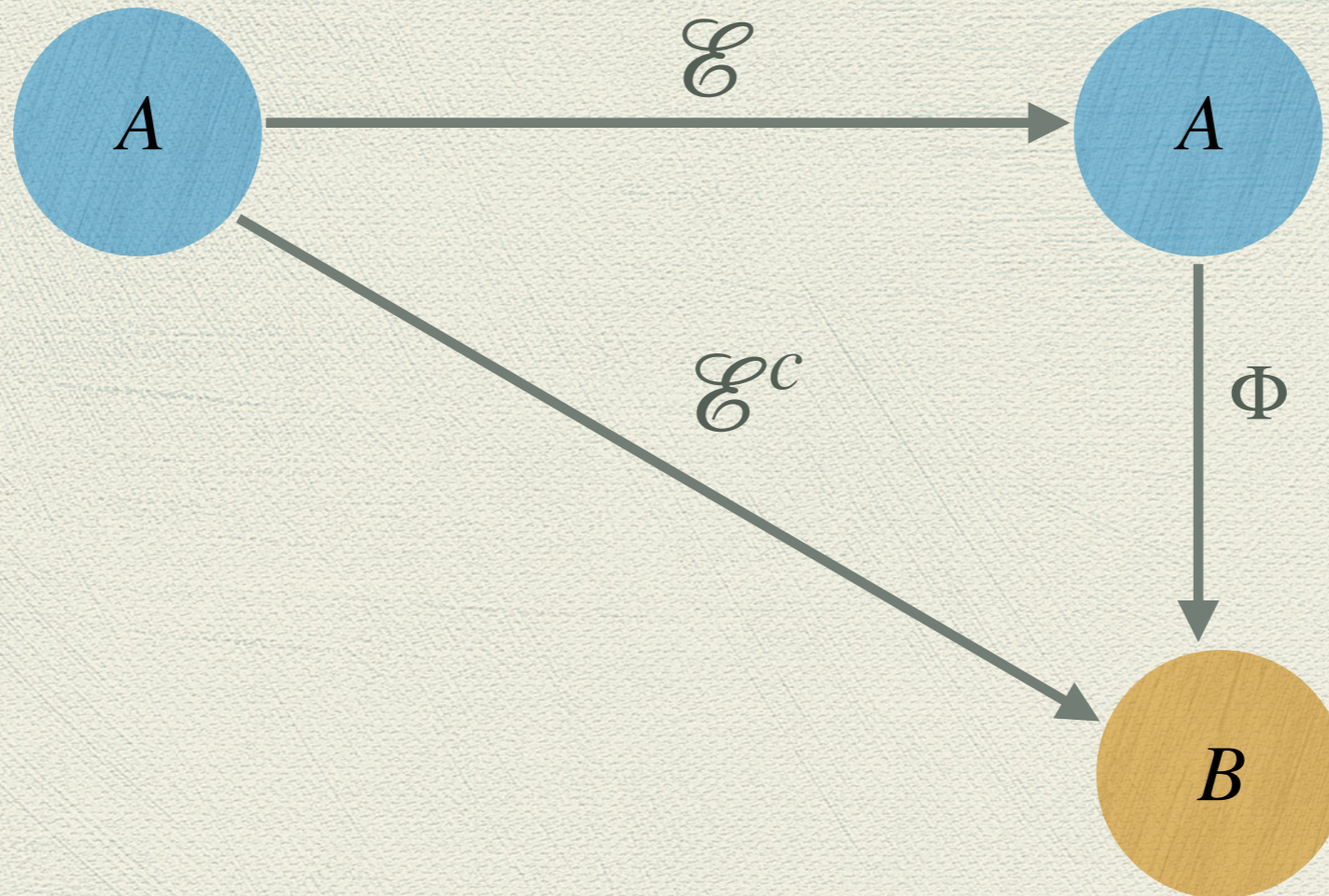
$$\mathcal{E}_U(\rho) = U\rho U^\dagger$$

$$\tilde{\mathcal{E}}_U(\rho) = \text{Tr}(\rho)$$

$$\mathcal{E}_\psi(\rho) = \text{Tr}(\rho) |\psi\rangle\langle\psi|$$

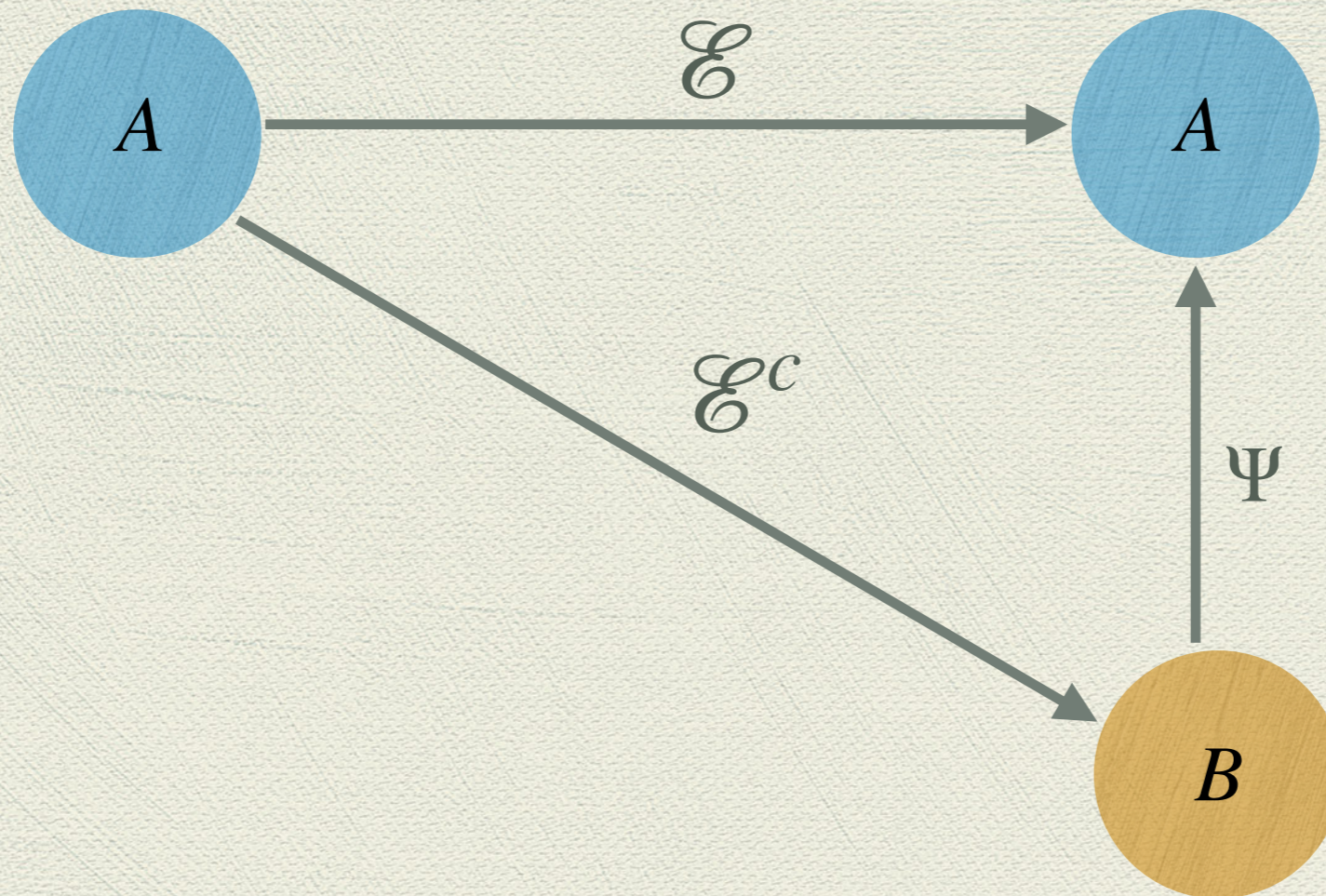
$$\mathcal{E}_\psi^c(\rho) = \rho$$

Degradable Channels



$$\mathcal{E}^c = \Phi \circ \mathcal{E}$$

Anti-Degradable Channels



$$\mathcal{E} = \Psi \circ \mathcal{E}^c$$

Examples of degradable and anti-degradable Channels

$$\mathcal{E}_0(\rho) = \text{Tr}(\rho) \frac{I}{d}$$

$$\tilde{\mathcal{E}}_0(\rho) = \frac{I}{d}(\rho \otimes I)$$

$$\Phi \circ \tilde{\mathcal{E}}_0 = \mathcal{E}_0$$

$$\Phi(X) = \text{Tr}_A(X)$$

Anti-Degradable

Examples of degradable and anti-degradable Channels

$$\mathcal{E}_U(\rho) = U\rho U^\dagger$$

$$\tilde{\mathcal{E}}_U(\rho) = \text{Tr}(\rho)$$

$$\tilde{\mathcal{E}}_U = \Phi \circ \mathcal{E}_U$$

$$\Phi(X) = \text{Tr}(X)$$

Degradable

Examples of degradable and anti-degradable Channels

$$\mathcal{E}_\psi(\rho) = \text{Tr}(\rho) |\psi\rangle\langle\psi| \quad \tilde{\mathcal{E}}_\psi(\rho) = \rho$$

$$\mathcal{E}_\psi = \Phi \circ \tilde{\mathcal{E}}_\psi$$

$$\Phi(X) = \text{Tr}(X) |\psi\rangle\langle\psi|$$

Anti-Degradable

Is there any non-trivial example of a degradable channel?

For which we can calculate the quantum capacity?

Q4: Private Capacity of Quantum Channels

For more information and for details, see:

[Capacities of the covariant Pauli channel](#)

A Poshtvan, V Karimipour

[Physical Review A 106 \(6\), 062408](#)

End of part III